

# anonymous routing and mix nets (Tor)

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Significant fraction of these slides are  
borrowed from CS155 at Stanford

# Anonymous web browsing

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## □ Why?

1. Discuss health issues or financial matters anonymously
2. Bypass Internet censorship in parts of the world
3. Conceal interaction with gambling sites
4. Law enforcement

## □ Two goals:

- Hide user identity from target web site: (1), (4)
- Hide browsing pattern from employer or ISP: (2), (3)

## □ Stronger goal: mutual anonymity (e.g. remailers)

# Current state of the world I

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- ❑ ISPs tracking customer browsing habits:
  - Sell information to advertisers
  - Embed targeted ads in web pages (1.3%)
    - » Example: MetroFi (free wireless)

[Web Tripwires: Reis et al. 2008]
- ❑ Several technologies used for tracking at ISP:
  - NebuAd, Phorm, Front Porch
  - Bring together advertisers, publishers, and ISPs
    - » At ISP: inject targeted ads into non-SSL pages
- ❑ Tracking technologies at enterprise networks:
  - Vontu (symantec), Tablus (RSA), Vericept

# Current state of the world II

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- EU directive 2006/24/EC: 3 year data retention
  - For ALL traffic, requires EU ISPs to record:
    - » Sufficient information to identify endpoints  
(both legal entities and natural persons)
    - » Session duration
      - … but not session contents
  - Make available to law enforcement
    - » … but penalties for transfer or other access to data
  
- For info on US privacy on the net:
  - “privacy on the line” by W. Diffie and S. Landau

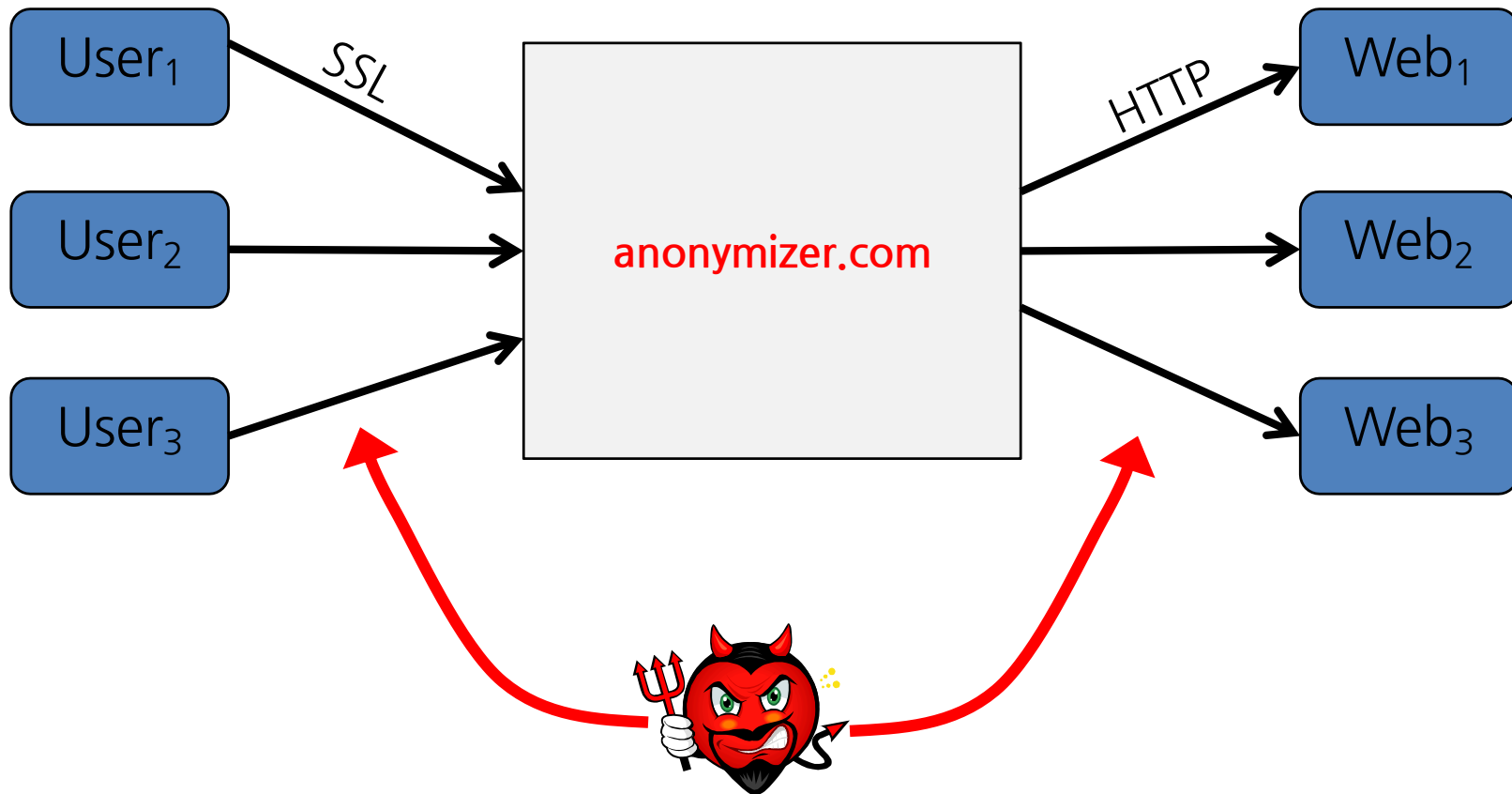
# Part 1: network-layer privacy

## Goals:

- Hide user's **IP address** from target web site
- Hide browsing destinations from network

# 1<sup>st</sup> attempt: anonymizing proxy

HTTPS:// anonymizer.com ? URL=target



# Anonymizing proxy: security

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- ❑ Monitoring ONE link: eavesdropper gets nothing
- ❑ Monitoring TWO links:
  - Eavesdropper can do traffic analysis
  - More difficult if lots of traffic through proxy

- ❑ Trust: proxy is a single point of failure
  - Can be corrupt or subpoenaed
    - » Example: The Church of Scientology vs. anon.penet.fi

- ❑ Protocol issues:
  - Long-lived cookies make connections to site linkable

# How proxy works

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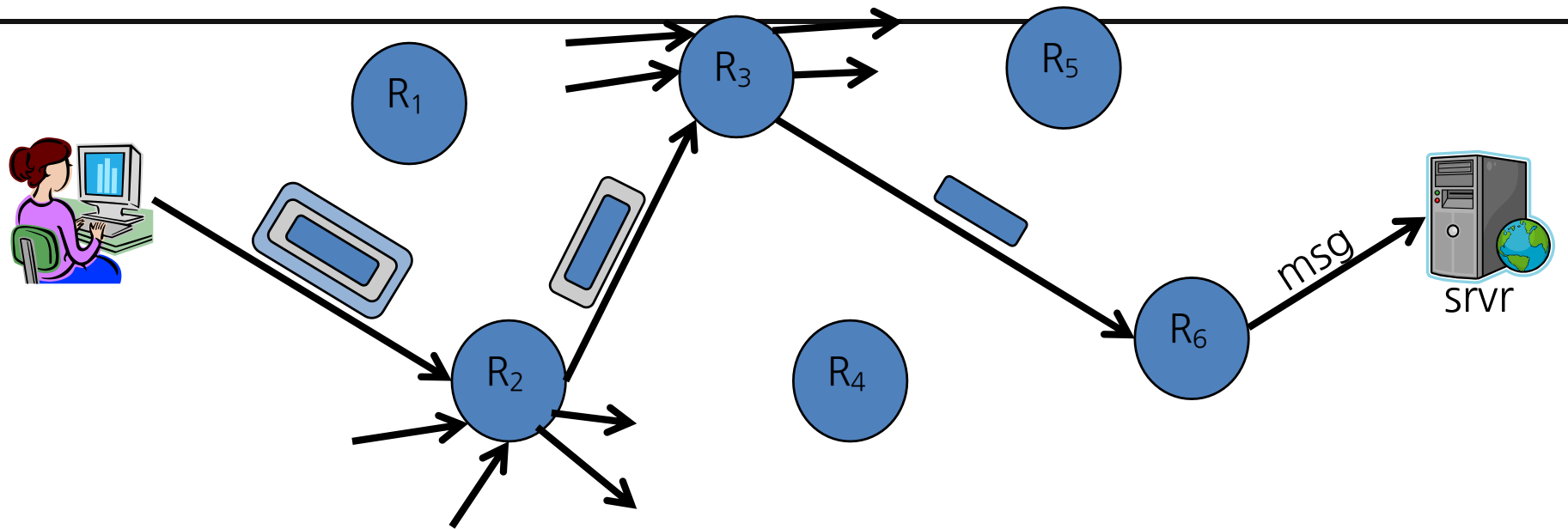
- ❑ Proxy rewrites all links in response from web site
  - Updated links point to anonymizer.com
    - » Ensures all subsequent clicks are anonymized
- ❑ Proxy rewrites/removes cookies and some HTTP headers
  
- ❑ Proxy IP address:
  - if a single address, could be blocked by site or ISP
  - anonymizer.com consists of >20,000 addresses
    - » Globally distributed, registered to multiple domains
    - » Note: chinese firewall blocks ALL anonymizer.com addresses
  
- ❑ Other issues: attacks (click fraud) through proxy



# 2<sup>nd</sup> Attempt: MIX nets

Goal: no single point of failure

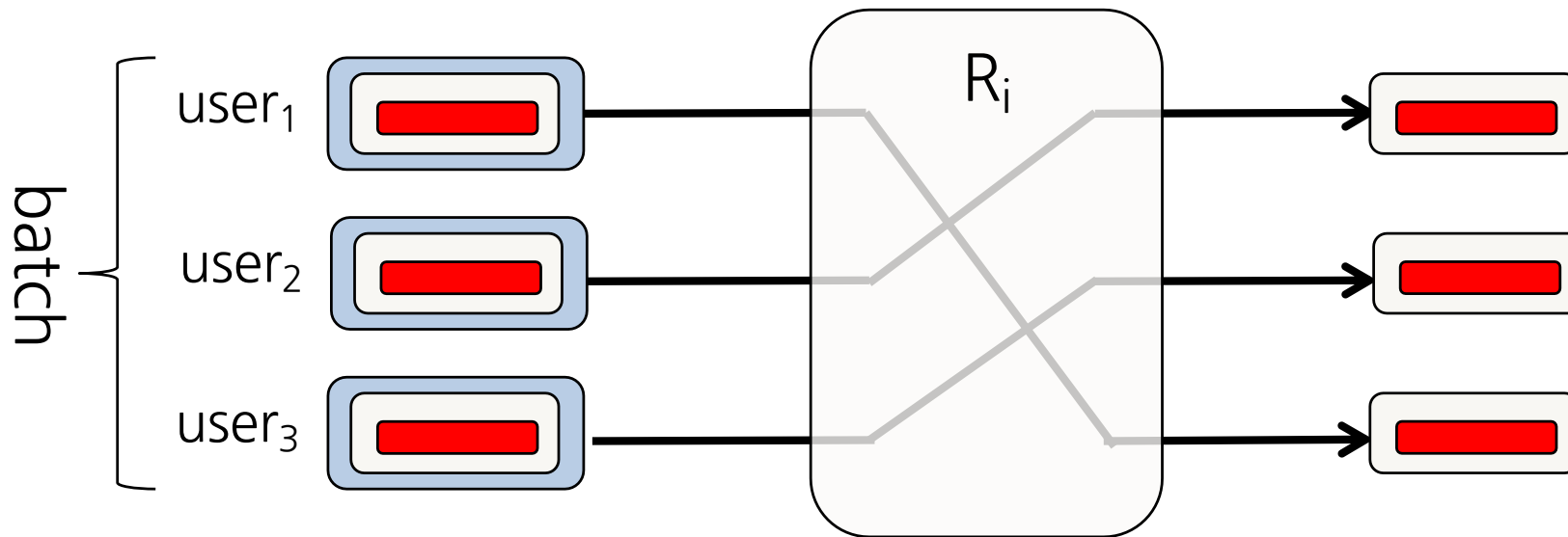
# MIX nets [Chaum'81]



- Every router has public/private key pair
  - Sender knows all public keys
- To send packet:
  - Pick random route:  $R_2 \rightarrow R_3 \rightarrow R_6 \rightarrow \text{srvr}$
  - Onion packet:

$E_{pk_2}(R_3, E_{pk_3}(R_6, E_{pk_6}(\text{srvr}, \text{msg})))$

# Eavesdropper's view at a single MIX



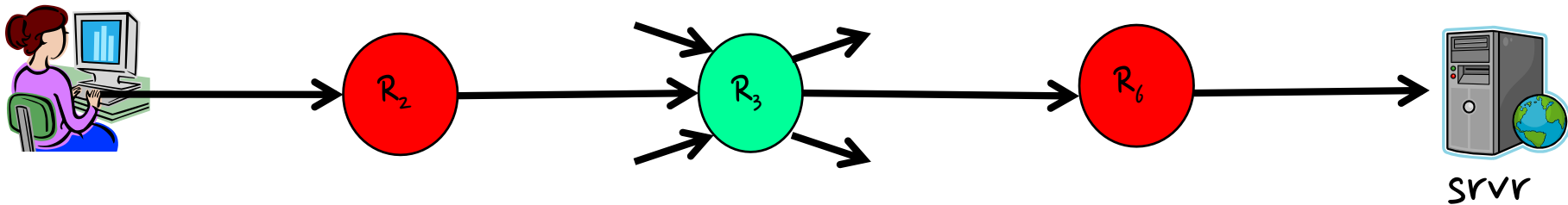
- ❑ Eavesdropper observes incoming and outgoing traffic
- ❑ Crypto prevents linking input/output pairs
  - Assuming enough packets in incoming batch
  - If variable length packets then must pad all to max len
- ❑ Note: router is stateless

# Performance

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- Main benefit:

- Privacy as long as at least one honest router on path



- Problems:

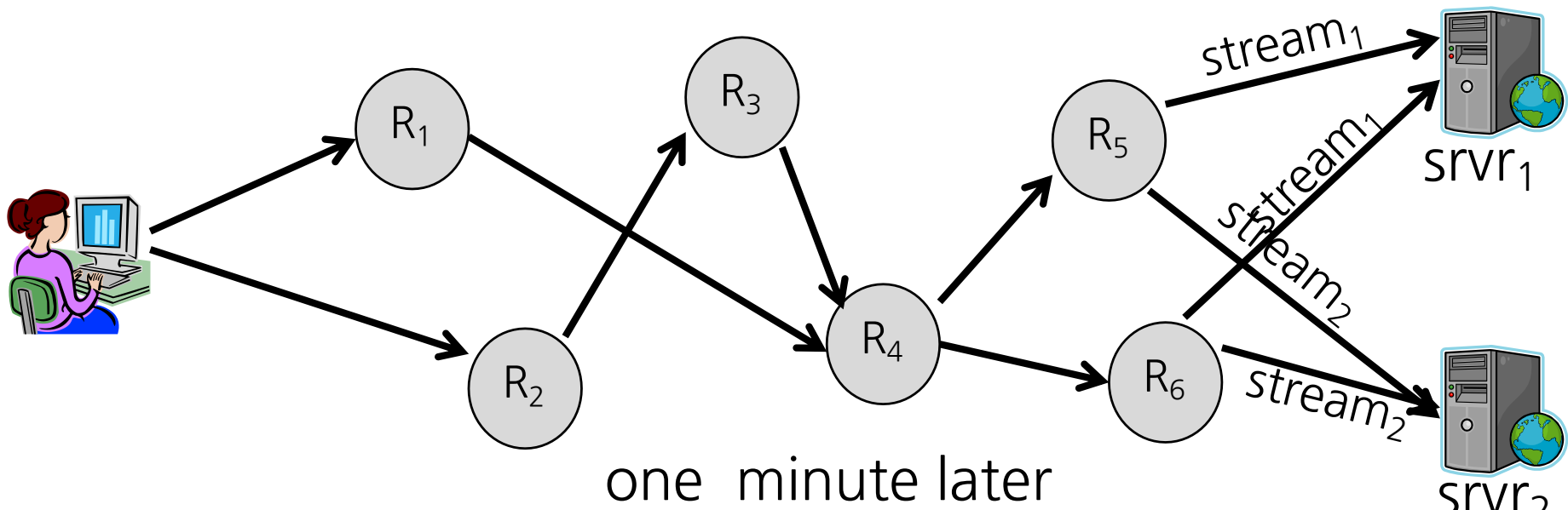
- High latency (lots of public key ops)
  - » Inappropriate for interactive sessions
  - » May be OK for email (e.g. Babel system)
- No forward security

# 3<sup>rd</sup> Attempt: Tor MIX circuit-based method

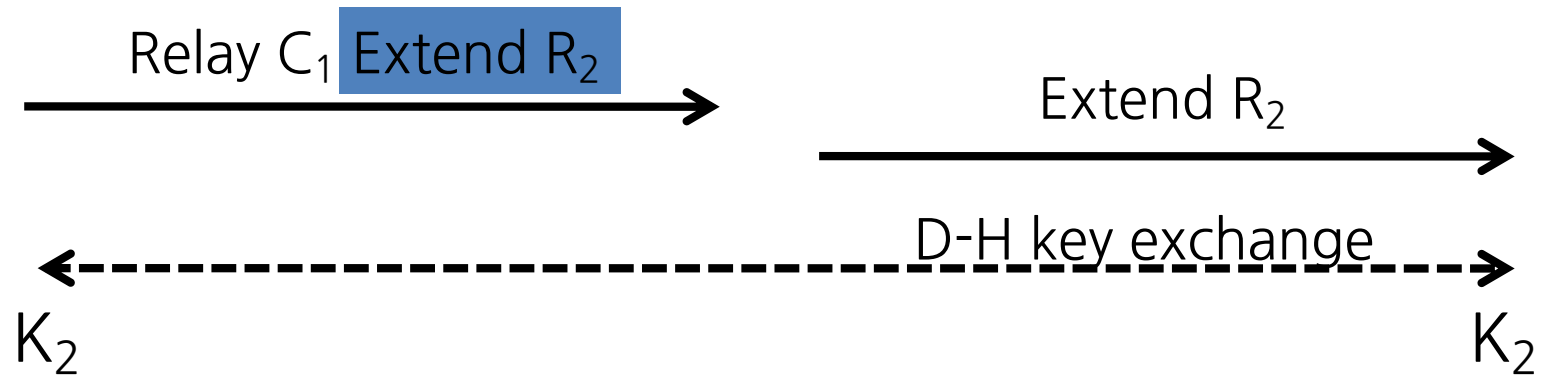
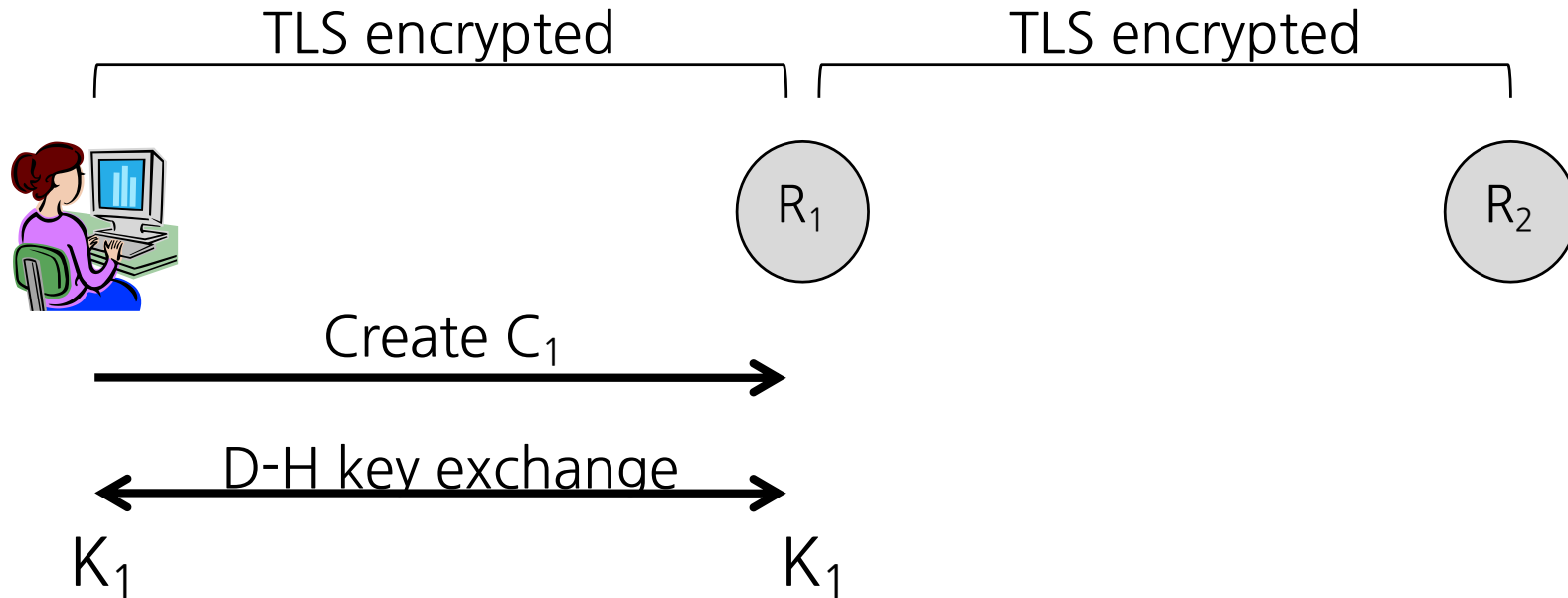
Goals: privacy as long as one honest  
router on path,  
and reasonable performance

# The Tor design

- ❑ Trusted directory contains list of Tor routers
- ❑ User's machine preemptively creates a circuit
  - Used for many TCP streams
  - New circuit is created once a minute

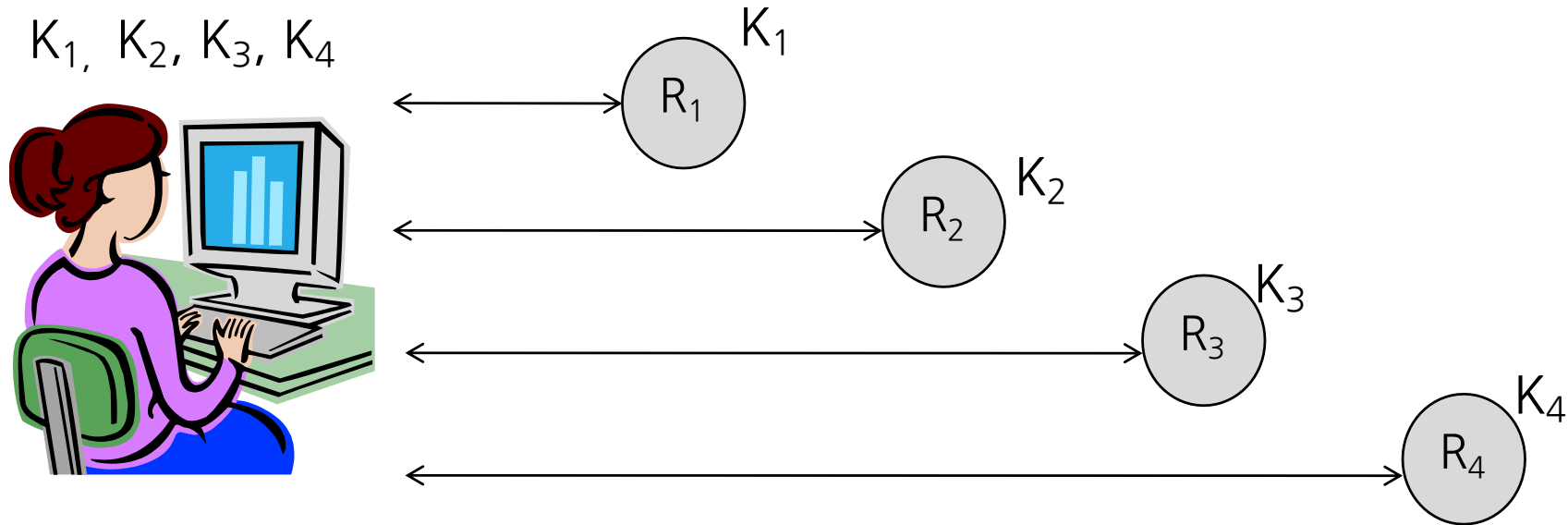


# Creating circuits



# Once circuit is created

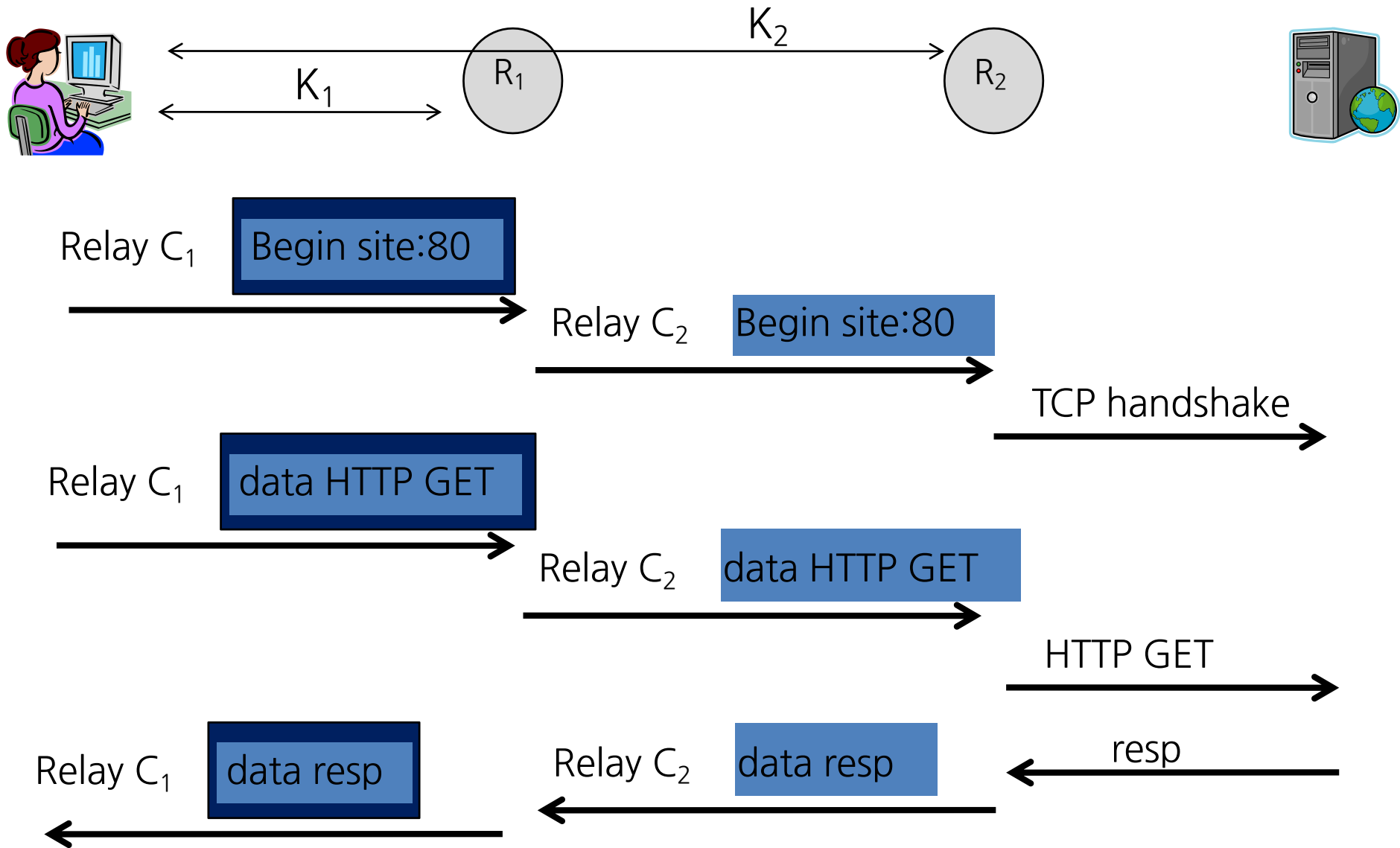
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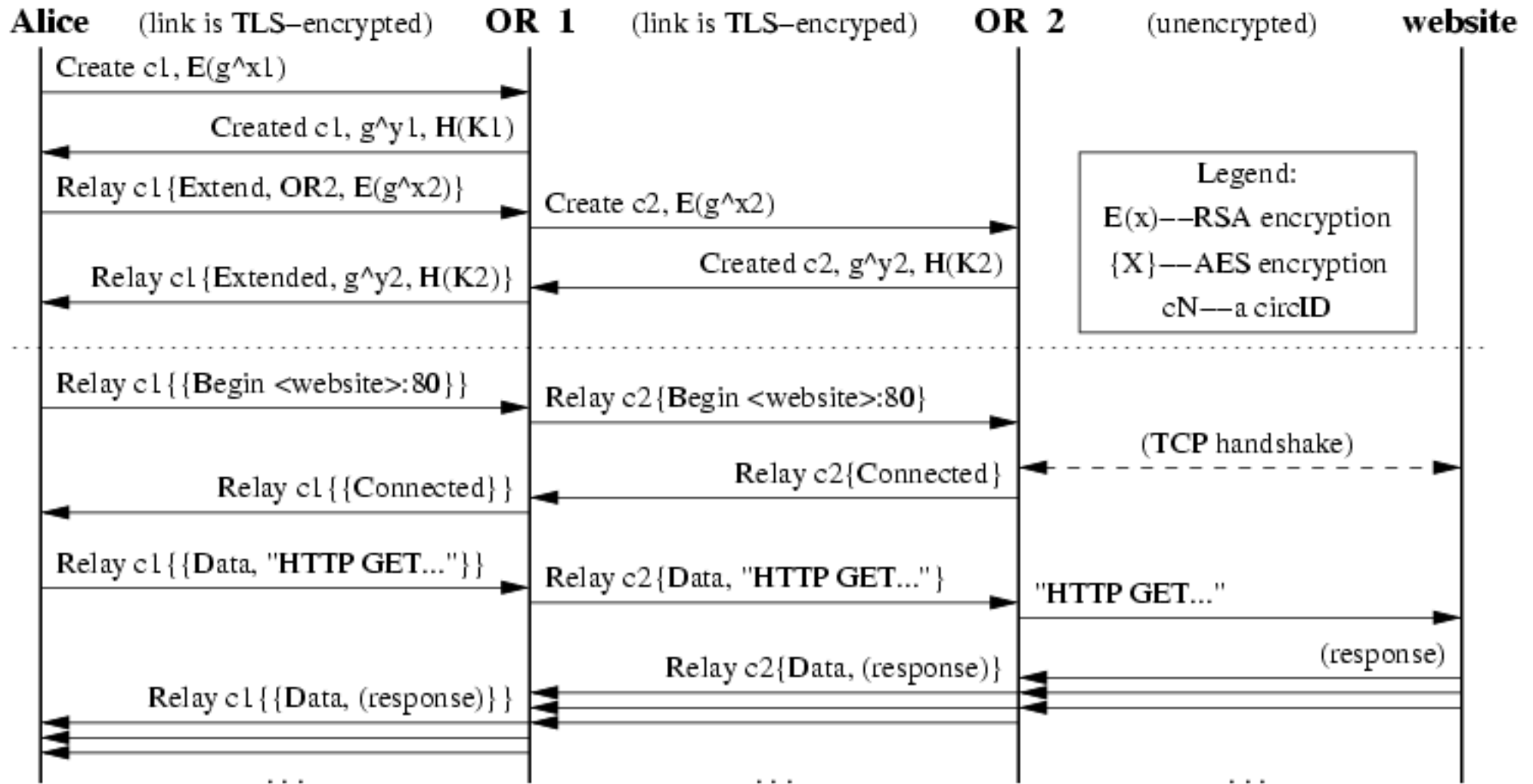
- ❑ User has shared key with each router in circuit
- ❑ Routers only know ID of successor and predecessor



# Sending Data



# Complete View



# Properties

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- ❑ Performance:
  - Fast connection time: circuit is pre-established
  - Traffic encrypted with AES: no pub-key on traffic
- ❑ Tor crypto:
  - provides end-to-end integrity for traffic
  - Forward secrecy via TLS
- ❑ Downside:
  - Routers must maintain state per circuit
  - Each router can link multiple streams via CircuitID
    - » all streams in one minute interval share same CircuitID

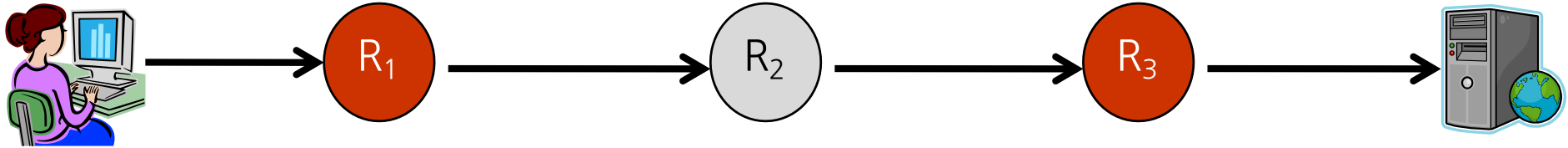
# Privoxy

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- ❑ Tor only provides network level privacy
  - No application-level privacy
    - » e.g. mail progs add “From: email-addr” to outgoing mail
  
- ❑ Privoxy:
  - Web proxy for browser-level privacy
  - Removes/modifies cookies
  - Other web page filtering

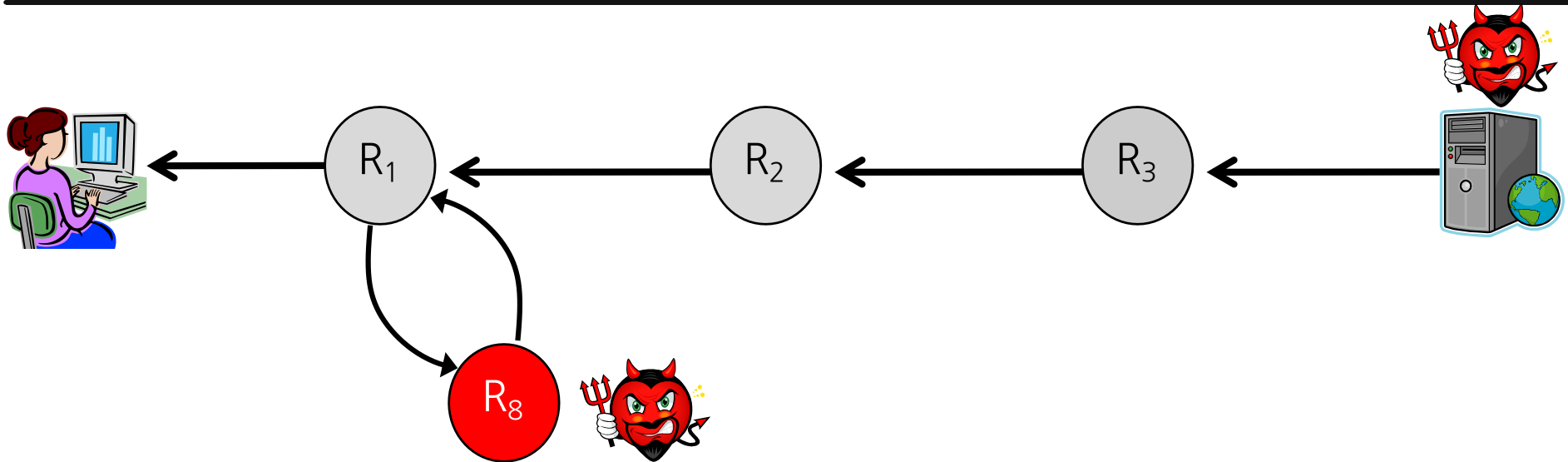
# Anonymity attacks: watermarking

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- Goal: R<sub>1</sub> and R<sub>3</sub> want to test if user is communicating with server
- Basic idea:
  - R<sub>1</sub> and R<sub>3</sub> share sequence:  $\Delta_1, \Delta_2, \dots, \Delta_n \in \{-10, \dots, 10\}$
  - R<sub>1</sub>: introduce inter-packet delay to packets leaving R<sub>1</sub> and bound for R<sub>2</sub>. Packet  $i$  delayed by  $\Delta_i$  (ms)
  - Detect signal at R<sub>3</sub>

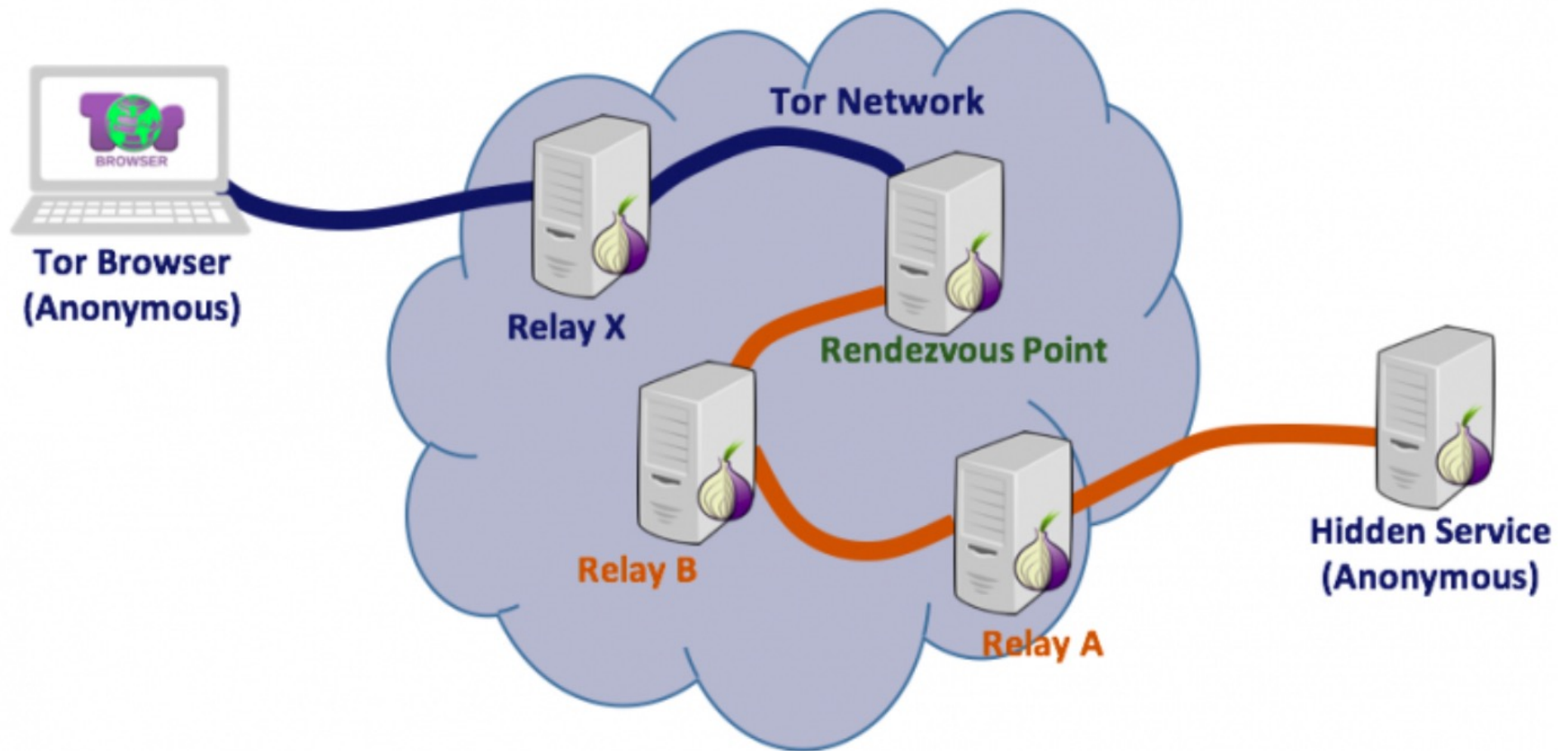
# Anonymity attacks: congestion



- ❑ Main idea:  $R_8$  can send Tor traffic to  $R_1$  and measure load on  $R_1$
- ❑ Exploit: malicious server wants to identify user
  - Server sends burst of packets to user every 10 seconds
  - $R_8$  identifies when bursts are received at  $R_1$   
Follow packets from  $R_1$  to discover user's ID

# Tor Hidden Service

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<https://www.torproject.org/docs/hidden-services.html.en>